

Extended Introduction to Computer Science

CS1001.py

Lecture 18:

Hash Tables (cont.)

Finite and Infinite Iterators

Instructors: Daniel Deutch, Amir Rubinstein
Teaching Assistants: Michal Kleinbort, Amir Gilad

Founding instructor: Benny Chor

School of Computer Science
Tel-Aviv University

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<http://tau-cs1001-py.wikidot.com>

Hash Tables - Reminder

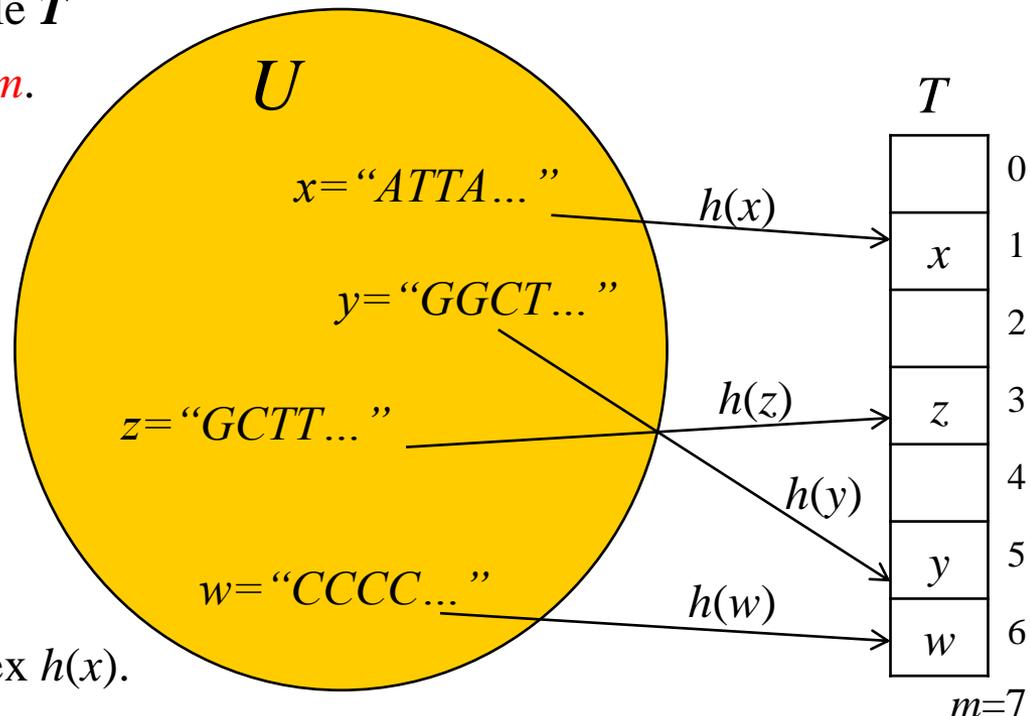
- Let us now define this mechanism more formally.
- Suppose elements belong to a large set (possibly infinite), called the "**universe**", denoted U (for example, all possible strings, all possible ID numbers etc.)
- Suppose we need to store n elements from U , and $n \ll |U|$.
(for example, genes of an organism, ID's of students in class right now)

- We can keep the elements in a table T called **hash table**, whose size is m .

$$|T| = m \ll |U|$$

- To map elements from U to T we will use a **hash function**

$$h: U \rightarrow \{0, 1, \dots, m-1\}$$



2 Element $x \in U$ will be stored at index $h(x)$.

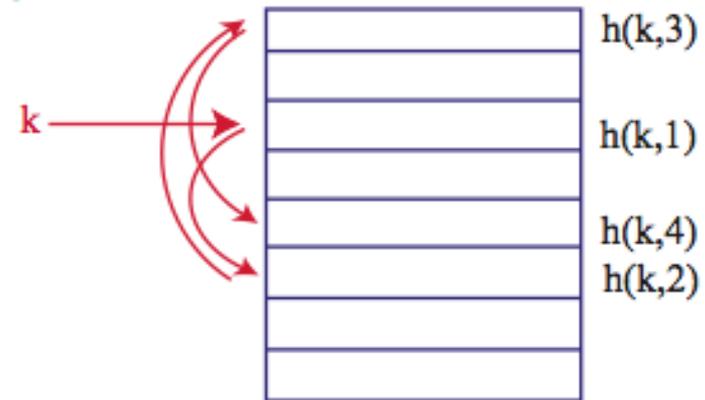
Two Approaches for Dealing with Collisions

- 1) **Chaining** – explained and implemented this ✓
- 2) **Open addressing** – we will briefly discuss it now

Two Approaches for Dealing with Collisions:

(2) Open Addressing

- In open addressing, each slot in the hash table contains **at most one** item. This obviously implies that **n cannot be larger than m**.
- Furthermore, an item will typically not stay statically in the slot where it "tried" to enter, or where it was placed initially. Instead, it may be moved a few times around.



- Open addressing is important in **hardware applications** where devices have many slots but each can only store one item (e.g. fast **switches** and high capacity **routers**). It is also used in **python dictionaries and sets**.
- There are many approaches to open addressing. We will describe a fairly recent one, termed **cuckoo hashing** (Pagh and Rodler, 2001).

Cuckoo Hashing: Motivation

- We saw that if $n \leq m$, hashing with chaining guarantees that insertion, deletion, and find are carried out in expected time $O(1)$ per operation, and **with high probability** (probability is over choices of inputs) $O(\log n / \log \log n)$ per operation. (The worst case time is $O(n)$ per operation.)
- In certain scenarios (e.g. fast routers in large internet nodes) we want **find** to run **with high probability** in $O(1)$ time. (The worst case time is still $O(n)$ per operation.)
- Compare $O(1)$ time **with high probability** to $O(1)$ **expected** time of hashing with chaining.
- **Cuckoo hashing** is one way to achieve this, but there are two prices to pay:
 - 1) Instead of $n \leq m$, we require $n \leq 7m/8$, or $n \leq 3m/4$, or $n \leq m/2$, or even $n \leq m/3$.
 - That is, we pay a price in terms of memory
 - 2) **insert** may take **somewhat longer time**.

Cuckoo Hashing

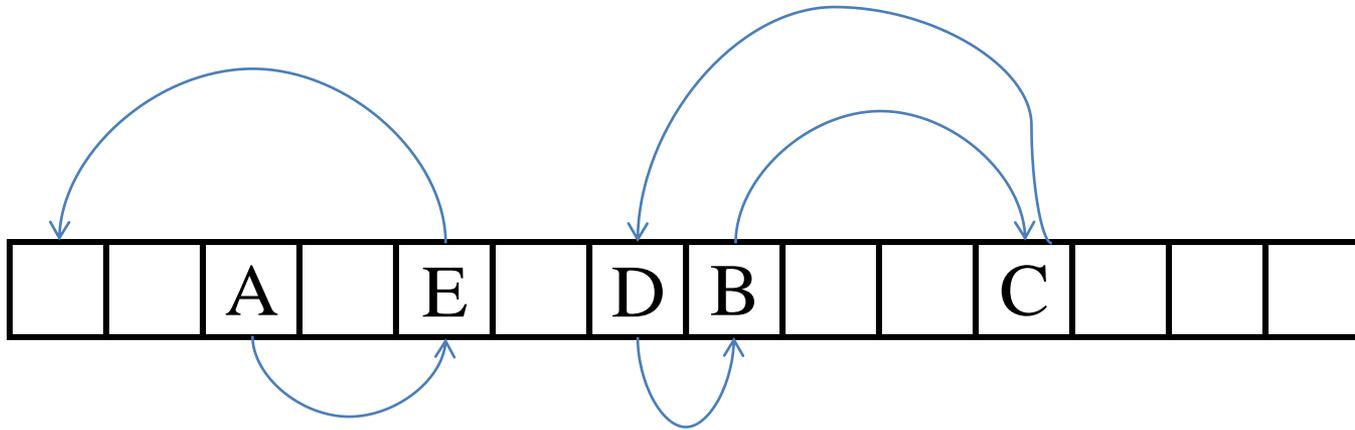
- **Cuckoo hashing** uses two distinct hash functions, h_1 and h_2 (improved versions use four, six, or eight, but the idea is the same).
- Each key, k , has **two potential slots** in the hash table, $h_1(k)$ and $h_2(k)$. If we search for k , all we have to do is look for it in these two locations (no chains here -- at most **one item** per slot).
- It is slightly more involved to **insert** a record whose key is k .

Cuckoo Hashing

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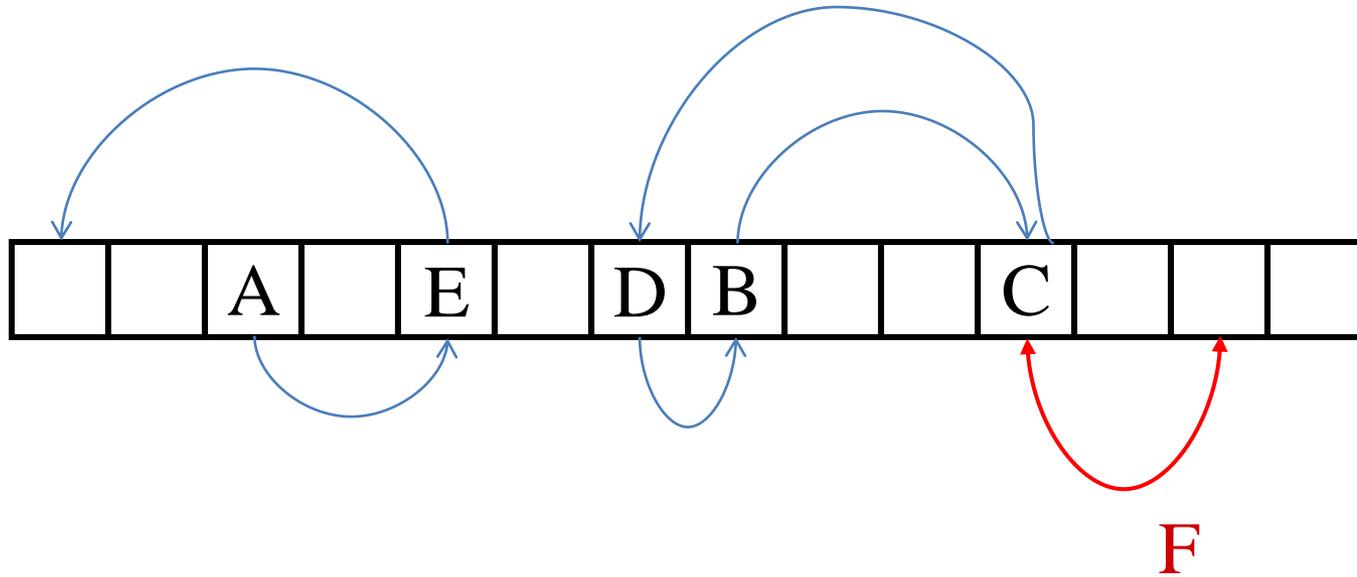
- If any of the two slots, $h_1(k)$ or $h_2(k)$ is empty, k is inserted there.
- If both slots are full, pick one of the two occupants, say x . Place k in x 's current slot.
- Assume this was location $h_1(x)$. Place x in its other slot, $h_2(x)$.
- If that slot was empty, we are done.
- Otherwise, the slot is occupied by some y . Place this y in its other slot, potentially kicking its present occupant, etc.,etc., until we find an empty slot.

Cuckoo Hashing: Examples



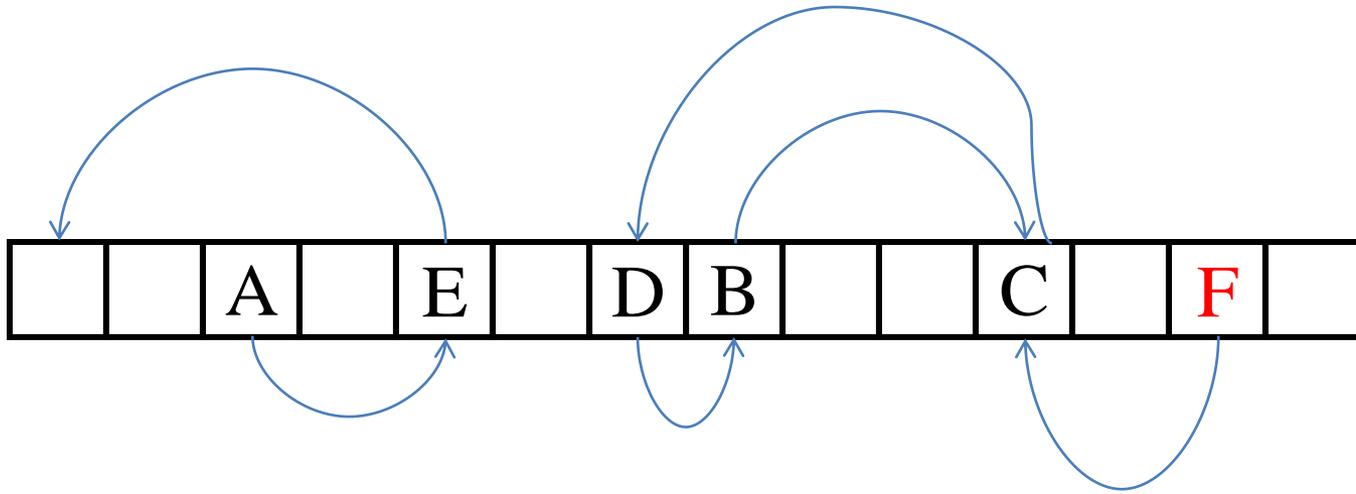
= The other potential slot for an item

Cuckoo Hashing: Examples



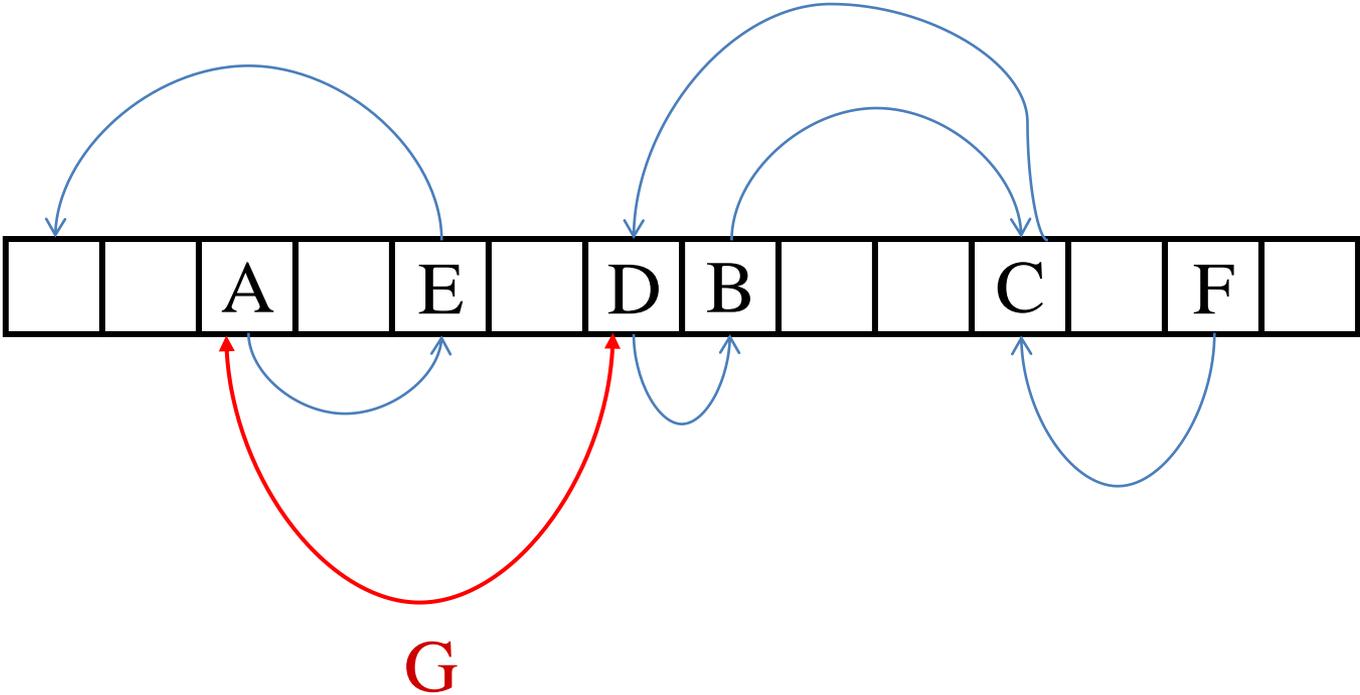
= The other potential slot for an item

Cuckoo Hashing: Examples



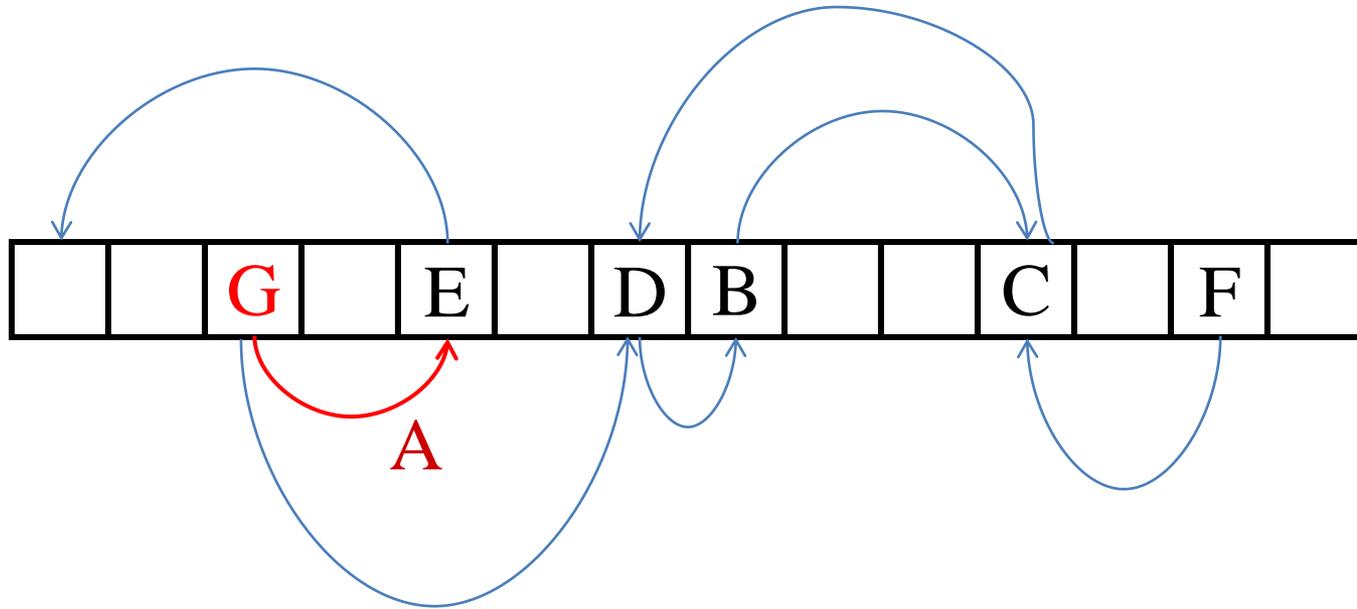
= The other potential slot for an item

Cuckoo Hashing: Examples



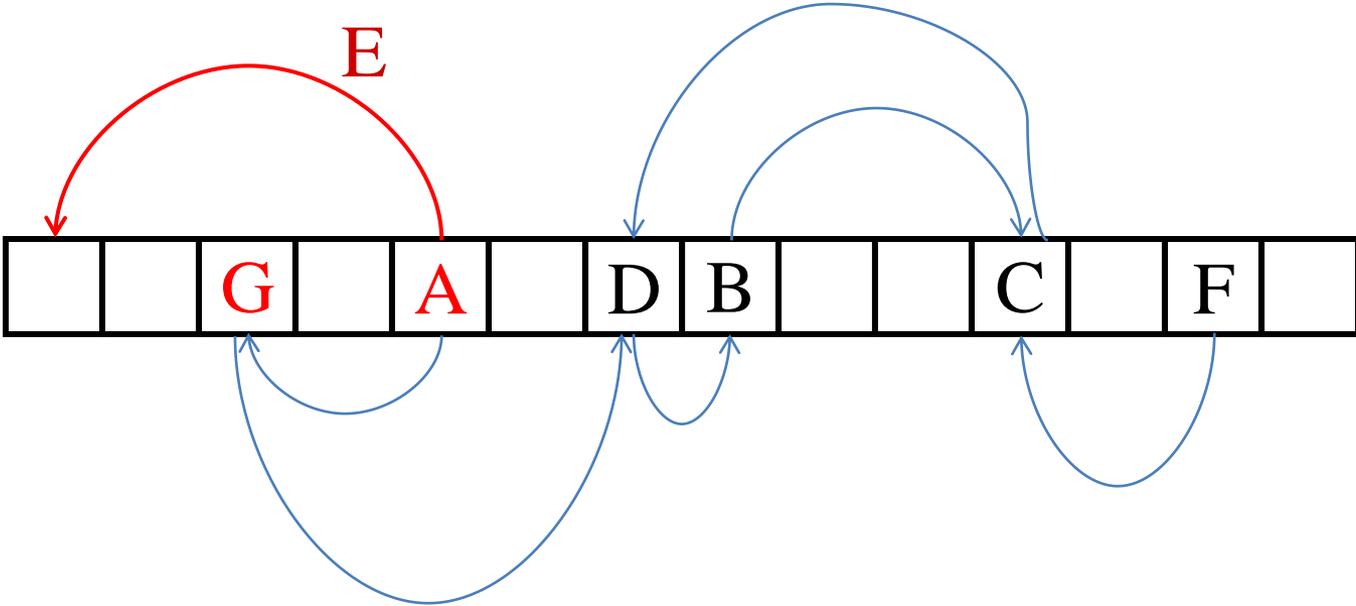
= The other potential slot for an item

Cuckoo Hashing: Examples



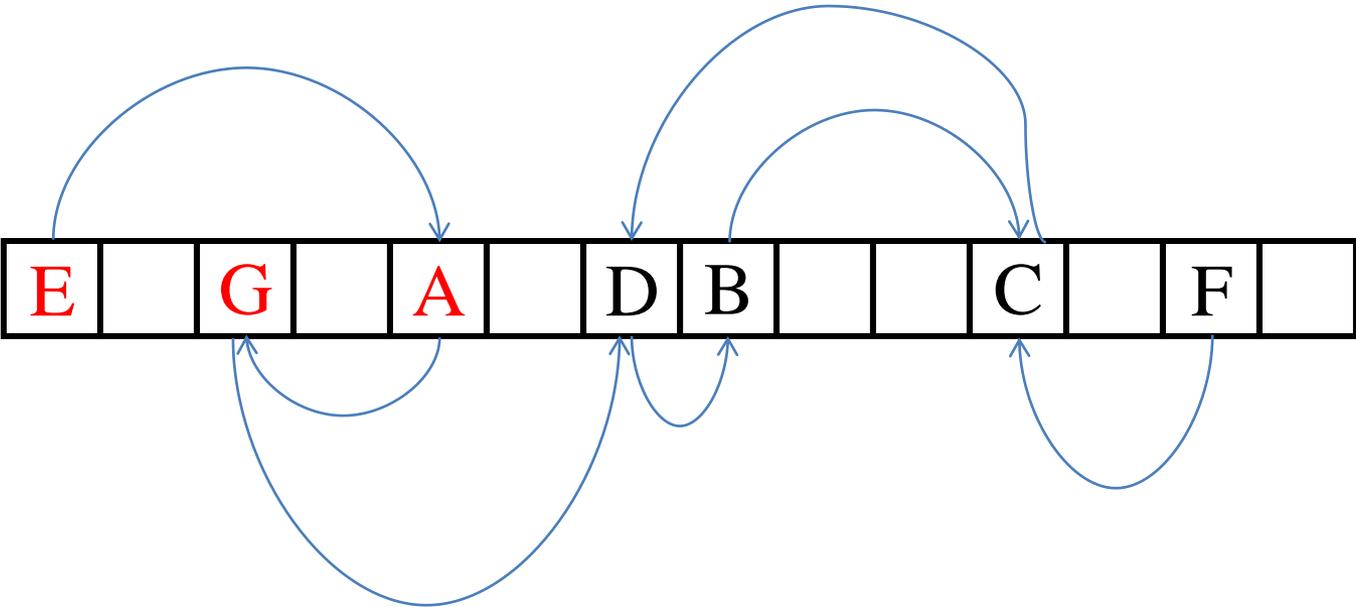
= The other potential slot for an item

Cuckoo Hashing: Examples



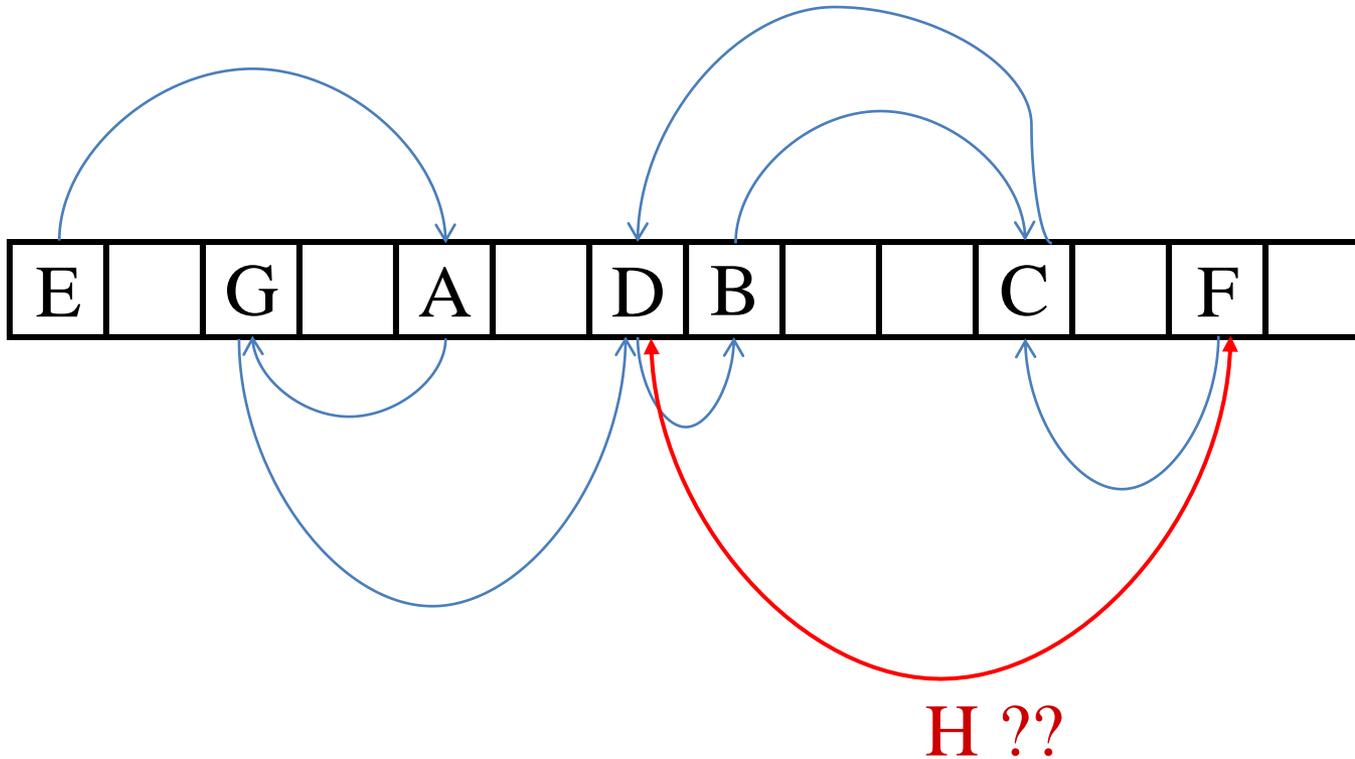
= The other potential slot for an item

Cuckoo Hashing: Examples



= The other potential slot for an item

Cuckoo Hashing: Examples



= The other potential slot for an item

Cuckoo Hashing - Deadlocks

- In the last example, we have reached a **cycle**, and we are in a non ending loop. This is called a **deadlock**.
- The union of the potential locations of **5 items** (B, C, D, F, H) is just **4 slots**.
- This obviously is very bad news for our cuckoo hashing.
- Notice that this is not a very likely event. With very high probability, the 10 potential locations ($10=5 \cdot 2$) will attain **more** than just 4 distinct values (which is why we got stuck in that the last example).

Cuckoo Hashing – Solving Deadlocks

- Another possible problem is that there will be no cycle, but the path leading to the successful insertions will be very **long**.
- Fortunately, such unfortunate cases occur with very low probability when the load factor , i.e. n/m , is **sufficiently low**. The common recommendation for two hash functions , $h_1(\cdot)$, $h_2(\cdot)$, is to have $n/m < 1/2$. (More hash functions enable a higher load factor).
- A theoretical solution: In case of failure (or very long path), **rehash** using **“fresh hash functions.”**
- A more practical solution: Maintain a very small **excess zone** (e.g. 32 excess slots for a hash table with $m=10000$ slots) and place items “causing trouble” there. If regular search (applying $h_1(x)$, $h_2(x)$) fails, search the excess zone as well.

Cuckoo Hashing in the Real World

- The load factor has to be smaller than 1. Yet a small load factor, say $n/m < 1/2$, is a **waste of memory**.
- In high performance routers, for example, most operations (including the hashing) are done **in silico, by the hardware**. The critical resource is memory area within the chip. Low load factor means wasted area.
- Instead of just 2 hash functions, **4 to 8** hash functions are utilized. This allows to increase the load factor to $n/m = 3/4$ or even $n/m = 7/8$.
- Suppose we use **4** hash functions, $h_1()$, $h_2()$, $h_3()$, $h_4()$. Given an element, x , that we wish to insert, we first check if any of the four locations $h_1(x)$, $h_2(x)$, $h_3(x)$, $h_4(x)$ is free.

Cuckoo Hashing in the Real World, cont.

- If these 4 locations are all taken, let a, b, c, d be the four elements in the above mentioned locations, respectively.
- Look, for example, at a . If one of the other 3 locations among $h_1(a), h_2(a), h_3(a), h_4(a)$ is free, we move a there, and put x in its place. If not, we do the same with respect to b , then c , then d .
- If all these are taken ($4+4\cdot 3=16$ different locations, typically), we go one more level down this search tree ($12\cdot 3 = 36$ additional locations, typically).
- If all these are taken, we give up on x and put it in the garbage bin (“excess zone” table).
- With very high probability, the small excess zone does not fill up. After removing elements from the table, we could try re-inserting such x to the hash table.

Designing Distinct Hash Functions

- Recall that the goal of designing a hash functions is that they map most sets of keys such that the maximal number of collisions is **small**.
- When having more than one hash functions, we have the additional goal that the different functions map same keys approximately **independently**. In Python, we could try variants of good ole hash.

For example:

```
def hash0(x):  
    return hash("0" + str(x))
```

```
def hash1(x):  
    return hash("1" + str(x))
```

```
def hash2(x):  
    return hash(str(x) + "2")
```

```
def hash3(x):  
    return hash(str(x) + "3")
```

Designing Distinct Hash Functions

A reminder concerning str (mapping objects to representing strings):

```
>>> [str(i) for i in range(10,20)]
['10', '11', '12', '13', '14', '15', '16', '17', '18', '19']
>>> str(2.2)
'2.2'
>>> str("2.2")
'2.2'
```

And now applying the four functions on a small domain:

```
>>> for f in (hash0,hash1,hash2,hash3):
    print([f(i) %23 for i in range(10,20)])
[ 2, 3, 4, 5, 6, 7, 8, 9, 10, 11]
[ 3, 2, 5, 4, 22, 21, 1, 0, 11, 10]
[12, 5, 17, 10, 16, 9, 7, 0, 10, 22]
[13, 4, 18, 9, 17, 8, 8, 22, 11, 21]
```

Random? Independent? Mixing well? You be the judges.

Cuckoo Hashing: Python Implementation and Analysis

Will not be done in class.

May appear in HW.

Hash Functions: Wrap Up

- Hash functions map large domains to smaller ranges.
- Example:
$$h : \{0,1,\dots,p^2\} \rightarrow \{0,1,\dots,p-1\},$$

defined by $h(x) = a \cdot x + b \pmod{p}$.
- Hash tables are extensively used for searching.
- If the range is larger than the domain, we cannot avoid **collisions** ($x \neq y$ with $h(x) = h(y)$). For example, in the example above, if $x_1 = x_2 \pmod{p}$ then $h(x_1) = h(x_2)$.
- If the range size is larger than **the square root** of the domain size, there will be **collisions** with high probability .
- A good hash function should create few collisions for most subsets of the domain ("few" is relative to size of subset).

Wrap Up: Dictionary in Computer Science

- In computer science, a **dictionary** is a data structure supporting efficient **insert**, **delete**, and **search** operations.
- This is an abstract notion, and should **not** be confused with Python's **class dict**, although **class dict** can be thought of as an implementation of the abstract data type **dictionary** (with elements that are pairs key:value).
- There are **two variations** of this data structure, according to the type of the elements stored in the data structure.
 - pairs **key: value**, or
 - Just **keys**

In any case, we assume all keys are **unique** (different items have different keys).

- Hash tables provide an excellent way to implement dictionaries.

Using Hash Functions and Tables: Wrap Up

- We explained chaining as a way to resolve collisions.
- We also studied the paradigm of cuckoo hashing , using two hash functions $h_1()$, $h_2()$ (or four, or eight).
Cuckoo hashing is aimed at a constant time find operation, with **high probability**, at the cost of a slightly longer insert operation.
- In the data structures course, you will see additional collisions resolution means, such as double hashing, etc.
- Python **sets** and **dictionaries** use hash tables, thus searching an element in a set / dict takes $O(1)$ time on average. Collisions are solved using open addressing, in a more sophisticated manner. In addition, the size of the hash table is dynamic.

Iterators and Generators

- Linked lists and Python's built-in lists (arrays) are two ways to represent a collection of elements. There are others, such as trees, dictionaries, and more.
- It is **desirable** that functions that use the data as part of a computation should be as oblivious as possible to such internal representation, which **may change over time** .
 - This general idea is captured in a concrete way by Python's **iterators**.
 - Iterators will provide a generic access to a collection of items. So generic that it will even allow us to access an **infinite** collection (also known as **stream**)!
 - Python's **generators** are tools to **create iterators** .

Iterables

- An **iterable** object is an object capable of returning its members one at a time.
- In particular, we can use a **for** loop on iterables
- Examples of iterables include:
 - all sequence types (such as list, str , tuple and range)
 - some non-sequence types like dict, set and File
 - and objects of any user defined classes with an `__iter__()` or `__getitem__()` method (that is, this is how you make your new class iterable. But we will not see this)

(see <http://docs.python.org/dev/glossary.html#term-iterable>)

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range is a special iterable class.

```
>>> a=range(10)
>>> type(a)
<class 'range'>
>>> a
range(0, 10)
>>> a[2]
```

Iterators

- An **iterator** is an object representing a stream of data.
- Each iterable type in Python has its own iterator type, created using the built-in `iter()`
- Repeated calls to the built-in function `next(it)`, where `it` is an iterator (or calls to the iterator's `__next__()` method) return successive items in the stream.
- When no more data are available a **StopIteration exception** is raised instead. At this point, the iterator object is "exhausted", and any further calls to `next(it)` just raise **StopIteration exception** again.

```
>>> it = iter([0,1,2])
```

```
>>> next(it)
```

```
0
```

```
>>> next(it)
```

```
1
```

```
>>> next(it)
```

```
2
```

```
>>> next(it)
```

```
Traceback (most recent call last):
```

```
  File "<pyshell#26>", line 1, in <module>
```

```
    next(it)
```

```
StopIteration
```

Iterables and Iterators

- We can create an iterator by calling the function `iter` with an **iterable object** argument (like list, tuple, str, dict, range , etc.)
- This function does not modify the original iterable object. In fact, when we loop over an iterable using **for**, an iterator is created first, and then the items are called, one by one, using `next()` .

```
>>> table = {"benny":72,"rani":82,"raanan":92}
```

```
>>> next(table)
```

```
Traceback (most recent call last):
```

```
  File "<pyshell#13>", line 1, in <module>
```

```
    next(table)
```

```
TypeError: dict object is not an iterator
```

```
>>> it = iter(table)
```

```
>>> next(it)
```

```
'rani'
```

```
>>> next(it)
```

```
33 'benny'
```

Iterables and Iterators, cont.

```
>>> next(it)
```

```
'raanan'
```

```
>>> next(it)
```

```
Traceback (most recent call last):
```

```
  File "<pyshell#18>", line 1, in <module>
```

```
    next(it)
```

```
StopIteration
```

```
>>> table = {"benny":72,"rani":82,"raanan":92}
```

```
>>> for key in table:           # an iterator is created
```

```
    print(key)                 # "under the hood"
```

```
                                # more details later
```

```
rani
```

```
benny
```

```
raanan
```

Iterables and Iterators, cont.

- As we see from this example, a dictionary (when transformed into an iterator), returns the keys one by one.
- Files return the lines one by one, etc.

- We can turn an iterator into a list as well. This list will reflect the current state of the iterator, **not** its original state:

```
>>> table = {"benny":72,"rani":82,"raanan":92}
```

```
>>> it = iter(table)
```

```
>>> next(it)
```

```
'rani'
```

```
>>> list(it)
```

```
['benny', 'raanan']
```

```
>>> next(it)
```

```
Traceback (most recent call last):
```

```
File "<pyshell#82>", line 1, in <module>
```

```
next(it)
```

```
StopIteration
```

Thou Shalt Not Modify an iterable during Iteration

- If we add or remove elements from an iterable during iteration, **strange things** may happen. For example

```
>>> elems = ['a','b','c']
>>> for e in elems:
    print(e)
    elems.remove(e)
```

a

c

```
>>> elems
['b']
>>>
```

adapted from

<http://unspecified.wordpress.com/2009/02/12/thou-shalt-not-modify-a-list-during-iteration/>

Iterators as a Tool for **Abstraction**

- The use of iterators hides the implementation of data collections. For example, when we see the code

for x **in** SomeCollection:

....

- We do not know if SomeCollection is a list, a dict, or any user defined data collection. Furthermore, we can later modify the implementation of SomeCollection, for example change it from a list to a dict, and the code using it (called **client code**) will not have to be changed.
- Similarly when we use next(**it**), **it** may be an iterator of any kind of a data collection, with any order of traversal.

Iterables, Iterators, and Generators: More Examples

```
>>> mylist = [x for x in range(10**8)]
>>> it1 = iter(mylist)
>>> it2 = (x for x in range(10**8)) #note the () instead of []

>>> type(mylist)
<class 'list'>
>>> type(it1)
<class 'list_iterator'>
>>> type(it2)
<class 'generator'> #list comprehension syntax inside () is one way to
                    #create a generator. The other will be shown next

>>> # mylist
    # typing this without the comment will clobber your screen
    # and most likely will cause your Python shell to crash

>>> it1
<list_iterator object at 0x17027d0>
>>> it2
<generator object <genexpr> at 0x1704f08>
```

Iterables, Iterators, and Generators, cont.

- it1 and it2 were iterators representing the first 10^8 integers.
- These integers can fit in just under 1GB RAM. But, can we have iterators representing even more items?

```
>>> it3 = (x for x in range(2**100)) #again, note the ()
```

```
>>> next(it3)
```

```
0
```

```
>>> next(it3)
```

```
1
```

- An iterable with 2^{100} elements will not fit in Amazon, Google, and NASA computers, even if taken together.
- Iterators and generators **represent streams**, but produce only **one element at a time**. Therefore, there is no problem representing a 2^{100} long stream!

Generators for Infinite Streams

In fact, there is no problem **representing streams** with **countably many** elements.

To do that, we will introduce generator functions.

So far, our functions contained no state, or memory. Successive calls to the function with the same arguments produced the same results (assuming the function does not refer to a global variable, which may have changed). This is now **going to change** .

```
def naturals():  
    """ a generator for all natural numbers """  
    n=1  
    while True:  
        yield n  
        n+=1
```

Generators for Infinite Streams, cont.

A function that contains a **yield** statement is termed a **generator function**. When a generator function is called, the **actual arguments** are bound to the function's **formal argument** names in the usual way, but no code in the body of the function is executed. Instead, a generator—iterator object is returned.

```
>>> naturals()
<generator object natural at 0x16f60d0>
>>> nat = naturals()
>>> nat
<generator object natural at 0x16f60a8
```

Generators, cont.

- nat is a generator--iterator. To get its "returned value", which is specified by the **yield** statement, we invoke next .

```
>>> next(nat)
```

```
1
```

```
>>> next(nat)
```

```
2
```

```
>>> [next(nat) for i in range(10)]
```

```
[3, 4, 5, 6, 7, 8, 9, 10, 11, 12]
```

- We see that nat has a **state** , which is retained, unchanged, between successive calls.
- We can have additional instances of the same generator function.

```
>>> nat2 = naturals()
```

```
>>> next(nat2)
```

```
1
```

```
>>> next(nat)
```

```
13
```

A Fibonacci Numbers Generator

```
def fib():  
    """ a generator for all Fibonacci numbers """  
    a, b = 0, 1  
    while True:  
        yield b  
        a, b = b, a+b
```

```
>>> Fib = fib()  
>>> Fib  
<generator object fib at 0x1704fa8>
```

- Again, Fib is a generator--iterator, so to get its "returned value", which is specified by the `yield` statement, we invoke `next()` .

```
>>> next(Fib)  
1  
>>> next(Fib)  
1  
>>> next(Fib)  
2  
>>> [next(Fib) for i in range(10)]  
[3, 5, 8, 13, 21, 34, 55, 89, 144, 233]
```

Execution Specification

When a **yield** statement

yield expression_list

is encountered, the state of the function is frozen , and the value of expression_list is returned to the caller of next.

By "frozen" we mean that all local state is retained, including the current bindings of local variables, the instruction pointer, and the internal evaluation stack: enough information is saved so that the next time next() is invoked, the function can proceed exactly as if the **yield** statement were just another external call.

(see <http://www.python.org/dev/peps/pep-0255/>)

This is similar to what is called **co-routine**, in contrast with a normal function call which is also called **sub-routine**.

Merging Sorted, Infinite Iterators

Suppose `iter1` and `iter2` are sorted iterators, and **both are infinite** . We wish to produce a new sorted iterator which is the merge of both.

```
def merge(iter1,iter2):  
    """ on input iter1, iter2, two infinite orted iterators,  
        produces the sorted merge of the two iterators """  
  
    left = next(iter1)  
    right = next(iter2)  
    while True:  
        if left<right:  
            yield left  
            left = next(iter1)  
        else:  
            yield right  
            right = next(iter2)
```

Merging Sorted, **Infinite** iterators: Execution

```
>>> nat1=natural()
```

```
>>> nat2=natural()
```

```
>>> nat3=merge(nat1,nat2)
```

nat3 , too is a generator--iterator, so to get its “returned value”, which is specified by the **yield** statement, we invoke next .

```
>>> next(nat3)
```

```
1
```

```
>>> next(nat3)
```

```
1
```

```
>>> next(nat3)
```

```
2
```

```
>>> next(nat3)
```

```
2
```

```
>>> [next(nat3) for i in range(10)]
```

```
[3, 3, 4, 4, 5, 5, 6, 6, 7, 7]
```

An Attempt to Merge Sorted, **Finite** iterators

Should the iterators in merge really be infinite ?

```
>>> nat1 = natural()
>>> nat2 = (n-2 for n in range(3))
>>> nat3 = merge(nat1,nat2)
>>> next(nat3)
-2
>>> next(nat3)
-1
>>> next(nat3)
0
>>> next(nat3)
```

Traceback (most recent call last):

File "<pyshell#48>", line 1, in <module>

next(nat3)

File "/Users/benny/Documents/IntroCS2011/Code/intro17/lecture17.py",

line 30, in merge

right=next(iter2)

StopIteration

What went wrong?

The merged iterator, **nat3**, was not yet exhausted, yet one of the arguments to merge, **nat2** was exhausted. The merging procedure still invoked `next(iter2)`. This has caused a `StopIteration` error.

Handling Errors: **try** and **except**

Python provides an elaborate mechanism to handle run time errors. For example, division by zero causes a `ZeroDivisionError` .

```
>>> 5/0
```

Traceback (most recent call last):

```
File "<pyshell#37>", line 1, in <module>
```

```
5/0
```

```
ZeroDivisionError: int division or modulo by zero
```

Such errors disrupt the flow of control in a program execution.

We may want to **detect** such error and allow the flow of control to **continue**. This may not be so important in the small programs written in this course, but becomes meaningful in large software projects.

Python enables such detection, using the keywords **try** and **except**.

Handling Errors: **try** and **except**: example

```
def division(a,b):  
    try:  
        return a/b  
    except ZeroDivisionError:  
        print("division by zero")
```

Let us now apply this function in two different cases:

```
>>> division(5,6)  
0.8333333333333334
```

```
>>> division(5,0)  
division by zero
```

We will employ this error handling mechanism to enable merging any non-empty sorted iterators, finite or infinite .

More on `try` and `except`

The example in the previous slide is not so good – we can solve this problem with an `if` statement. The following example shows a situation where we would need to write many `if` statements, so `try/except` is better

```
def compute(...):  
    try  
        # a long computation, with several steps  
        # that may cause zero divide  
    except ZeroDivisionError:  
        print("division by zero")
```

We will also use `try/except` when it is either `impossible` or `expensive` to check for the condition in advance. Example – when we invert a matrix, checking in advance that it is not singular would take as much time as inverting, so it makes more sense to try to invert, and raise an `exception` if we discover that the matrix is singular while we do it.

We can have multiple `except` clauses; a list of exceptions to be handled in each clause; and the last clause may omit exception names (to handle all others)

For loop

We mentioned that a **for** loop over an iterable using **for**, actually uses an iterator. We now show an example:

```
>>> elems = ['a','b','c']
```

```
for e in elems:  
    print(e)
```

```
a  
b  
c
```

Is the same as

```
It = iter(elems)  
while True:  
    try:  
        print(next(it))  
    except StopIteration:  
        break
```

```
a  
b  
c
```

Merging **Any** Non-Empty, Sorted iterators

```
def merge3(iter1,iter2):
```

```
    """ on input iter1, iter2, two non-empty sorted iterators, not  
    necessarily infinite, produces sorted merge of the two iterators """
```

```
    left=next(iter1)
```

```
    right=next(iter2)
```

```
    while True:
```

```
        if left<right:
```

```
            yield left
```

```
            try:
```

```
                left=next(iter1)
```

```
            except StopIteration:          # iter1 is exhausted
```

```
                yield right
```

```
                remaining = iter2
```

```
                break
```

merge3 : cont.

else:

yield right

try:

right=next(iter2)

except StopIteration: # iter2 is exhausted

yield left

remaining = iter1

break

end of the while loop

for elem in remaining: # protects against StopIteration

yield(elem)

Merge3: Examples of Executions

```
>>> iter1=(x**2 for x in range(4))
>>> iter2=natural()
>>> merged=merge3(iter1,iter2)
>>> [next(merged) for i in range(14)]
[0, 1, 1, 2, 3, 4, 4, 5, 6, 7, 8, 9, 9, 10]
```

```
>>> iter1=(x**2 for x in range(5))
>>> iter2=(x**3 for x in range(6))
>>> merged=merge3(iter1,iter2)
>>> [next(merged) for i in range(11)]
[0, 0, 1, 1, 4, 8, 9, 16, 27, 64, 125]
```

Finally, lets see what happens when the original iterators/generators are not sorted .

```
>>> iter1=(-1)**x*x**2 for x in range(5))
>>> iter2=(x**3 for x in range(6))
>>> merged=merge3(iter1,iter2)
>>> [next(merged) for i in range(11)]
[0, 0, -1, 1, 4, -9, 8, 16, 27, 64, 125]
# garbage in, garbage out
```